

Distributed Systems: Concepts and Design *Edition 3*

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Errata List

As with virtually all books, some bugs were discovered after printing. These errors will be corrected in subsequent printings (impressions). The corrections recorded to date are listed below.

First and second impressions

| Page | existing version | correction | date |
|------|------------------|--|----------|
| 33 | line 8 | replace "Common" by "Component" | 12-9-00 |
| 49 | line -2 | replace first sentence by "The delay between the start of a message's transmission from one process and the beginning of its receipt by another is referred to as <i>latency</i> " | 19-2-01 |
| 66 | lines 12 to14 | replace "The Internet is constructed from many <i>subnets</i> employing a variety of network technologies. A subnet is a set of interconnected nodes, all of which employ the same technology to communicate amongst themselves." with "The Internet is constructed from many <i>subnets</i> . A subnet is a unit of routing (delivering data from one part of the Internet to another); it is a collection of nodes that can all be reached on the same physical network." | 21-02-01 |
| 67 | lines 16 to18 | replace with "Latency is the delay that occurs after a send operation is executed before data starts to arrive at the destination computer. It can be measured as the time required to transfer an empty message. Here we are considering only network latency, which forms a part of the process-to-process latency defined in Section 2.3.1." | 21-2-01 |
| 68 | line -1 | delete "timing" | 19-2-01 |

| 93 | Figure 3.15 | Class E addresses need no have a 0 in the fifth binary digit. Delete the box containing a 0 in the row for Class E addresses and replace 27 by 28 in the arrow label. | 15-03-02 |
|-----|-------------|---|----------|
| | | Class E (reserved): 1 1 1 1 unused | |
| 94 | Figure 3.16 | bottom right cell, replace "128" by "240" and replace "247" by "255" | 15-01-01 |
| 94 | line 5 | replace "set to 1" by "set to 0" | 26-02-03 |
| 96 | lines 8-12 | replace with: | 19-2-01 |
| | | correspondence between IP addresses and addresses for the hosts on the local network and is done using an address resolution protocol (ARP). | |
| | | We now outline the implementation of an ARP for Ethernets. It uses dynamic enquiries in order to operate correctly when computers are added to a local network but exploits caching to minimize enquiry messages. | |
| 98 | line 16 | replace "195.203.255.255" by "203.255.255.255" | 9-10-00 |
| 108 | line 4 | replace "model" by "modem" | 20-2-03 |
| 128 | line 29 | replace "receiver" by "receiver (multicast ports are an exception, see Section 4.5.1) | 19-2-01 |
| 130 | line 19 | replace "Any process needing to send or receive messages must first create a socket bound to an Internet address and local port" by "To send or receive messages a process must first create a socket bound to an Internet address of the local host and a local port." | 19-2-01 |
| 130 | line -10 | replace first sentence by "Sockets normally provide non-blocking <i>sends</i> and blocking <i>receives</i> for datagram communication(a non-blocking <i>receive</i> is an option in some implementations)." | 19-2-01 |
| 132 | Figure 4.3 | aSocket.close has been moved to a <i>finally</i> clause at the end of the <i>main</i> function. The corrected version of this program is available in the file <i>UDPClient.java</i> on this website under Additional Material for Chapter 4. | 19-2-01 |
| 133 | Figure 4.4 | aSocket.close has been moved to a finally clause at the end of the main function. The corrected version of this program is available in the file UDPServer.java on this website under Additional Material for Chapter 4. | 19-2-01 |
| 136 | Figure 4.5 | aSocket.close has been moved to a finally clause at the end of the main function. The corrected version of this program is available in the file TCPClient.java on this website under Additional Material for Chapter 4. | 19-2-01 |
| 137 | Figure 4.6 | aSocket.close has been moved to a finally clause at the end of the main function. The corrected version of this program is available in the file TCPServer.java on this website under Additional Material for Chapter 4. | 19-2-01 |

| 152 | Figure 4.15 | replace "URL" by "URL or pathname" | 18-9-00 |
|-----|-------------|---|----------|
| 152 | last line | add "Proxies need the whole URL as shown in the figure, but it is advantageous to send only the pathname in the case of an origin server because there is less to send." | 18-9-00 |
| 152 | Figure 4.15 | replace "//www.dcs.qmw.ac.uk/index.html" by "http://www.dcs.qmw.ac.uk/index.html" | 2-3-02 |
| 154 | line -10 | replace "Figure 3.16" by "Figure 3.15" | 13-11-00 |
| 156 | Figure 4.17 | aSocket.close has been moved to a finally clause at the end of the main function. The corrected version of this program is available in the file MulticastPeer.java on this website under Additional Material for Chapter 4. | 19-2-01 |
| 169 | line 2 | replace "Section 4.< <cdr>" by "Section 4.3.1"</cdr> | 15-01-01 |
| 169 | line 12 | replace "Common" by "Component" | 12-9-00 |
| 205 | last line | replace "Exercise 5.13" by "Exercise 5.15" | 27-11-01 |
| 218 | line 24 | replace first two sentences of paragraph 4 by: "Migratory load-sharing systems can shift load at any time, not just when a new process is created. They use a mechanism called process migration: the transfer of an executing process from one node to another." | 15-10-01 |
| 274 | line 5 | replace "idempotent" by "self-inverse" | 17-01-02 |
| 291 | line -10 | replace "[Steiner et al. 1988]" by "[Neumann and T'so 1994]" | 21-9-00 |
| 293 | line -6 | Insert after "Kerberos Version 5" "[Neumann and T'so 1994] | 15-9-00 |
| 321 | line -6 | replace "There are three operations for altering directories: <i>AddName</i> , <i>ReName</i> and <i>UnName</i> ." by "There are two operations for altering directories: <i>AddName</i> and <i>UnName</i> ." | 8-11-00 |
| 326 | Figure 8.9 | replace $link(newdirfh, newname, dirfh, name) \rightarrow status$ Creates an entry $newname$ in the directory $newdirfh$ which refers to file $name$ in the directory $dirfh$. with $link(newdirfh, newname, fh) \rightarrow status$ Creates an entry $newname$ in the directory $newdirfh$ which refers to the file or directory fh . | 02-04-03 |

| 328 | lines 14-16 | replace The file handle returned in the previous step is used as a parameter in the next <i>lookup</i> step; the file system identifier in the file handle is first compared with the entries in the remote mount table held in the client to see whether another remote-mounted file store should be accessed. with The file handle returned in the previous step is used as a parameter in the next lookup step. Since file handles are opaque to NFS client code, the virtual file system is responsible for resolving file handles to a local or a remote directory and performing the necessary indirection when it references a local mount point. | 09-04-03 |
|-----|-----------------------------------|---|----------|
| 330 | line -8 | insert at start of paragraph There is one value of Tm_{server} for all the data blocks in a file and another for the file attributes. | 09-04-03 |
| 341 | line 9 | replace "For each file with a valid token, Venus must send a cache validation request" by "Before the first use of each cached file or directory after a restart, Venus therefore generates a cache validation request " | 8-11-01 |
| 366 | line 10 | replace "IP" by "IN" | 19-4-01 |
| 387 | line 16 | replace "17" by "16" | 8-11-01 |
| 417 | Exercise 10.11 | replace "Vi[j]" by "Vj[i]" | 12-9-00 |
| 440 | penultimate line | replace "protocol" by "description" | 14-9-00 |
| 441 | last paragraph, sentences 2 and 3 | replace by: "The validity property holds because IP multicast has that property. For agreement we require, first, that a process can always detect missing messages. That in turn means that it will always receive a further message that enables it to detect the omission. As this simplified protocol stands, we guarantee detection of missing messages only in the case where correct processes multicast messages indefinitely." | 19-10-00 |
| 441 | last paragraph sentence 4 | replace "The agreement property holds as long as" by " Second, the agreement property requires that" | 12-10-00 |
| 442 | line 1 | replace first sentence by: "Neither of the assumptions we made to ensure agreement is practical (see Exercise 11.14)." | 12-10-00 |
| 442 | line 2 | replace "However, validity and agreement are practically addressed in the protocols" by: "However, agreement is practically addressed in the protocols | 12-10-00 |

| 446 | Figure 11.14 | replace $On \ B-deliver(<"order",\ i,\ S>)\ with\ g=group(m)$ wait until $< m,\ i>$ in hold-back queue and $S=r_g+1;$ TO -deliver $m;$ //(after deleting it from the hold-back queue) $r_g=S;$ | 20-03-01 |
|-----|----------------------|--|----------|
| | | by On B-deliver($m_{order} = <$ "order", i , $S>$) with $g = group(m_{order})$ wait until $<$ m, $i>$ in hold-back queue and $S=r_g$; TO -deliver m ; //after deleting it from the hold-back queue) $r_g = S+1$ | |
| 449 | Figure 11.16 | replace "On B-deliver($$) from p_j , with $g=group(m)$ " by "On B-deliver($$) from $p_i(j\neq i)$, with $g=$ | 10-04-03 |
| | | group(m)" | |
| 454 | line -10 | Insert "For the case where a majority of processes are correct," before "we construct a solution to C from IC'. | 23-07-02 |
| 455 | line 4 | Insert "In systems with crash failures," before beginning of paragraph "Solving consensus is equivalent". | 23-07-02 |
| 457 | line 3 | replace "Figure 11.18" by "Figure 11.19" | 27-10-02 |
| 463 | Exercise 11.13 | replace first sentence by "Explain whether the algorithm for reliable multicast over IP works for open as well as closed groups" | 14-9-00 |
| 482 | line 9 | replace "Transfer \$100 from A to B" by "Transfer \$100 from B to A" | 16-8-02 |
| 498 | line 22 | replace " $T_{ u}$ " by " T_{j} " | 12-11-01 |
| 551 | Exercise 13.9 line 3 | replace "Figure 13.6" by "Exercise 13.4" | 13-7-00 |
| 562 | lines -5 to -11 | replace paragraphs on Agreement and Integrity with: | 26-2-03 |
| | | Agreement: Correct processes deliver the same sequence of views (starting from the view in which they join the group), and the same set of messages in any given view. That is, if a correct process delivers message m in view $v(g)$, then all other correct processes that deliver m also do so in the view $v(g)$. | |
| | | Integrity: If a correct process p delivers message m , then it will not deliver m again. Furthermore, $p \in group(m)$ and the process that sent m is in the view in which p delivers m . | |
| 639 | Figure 16.2 line 2 | replace "inta" with "int a" | 14-9-00 |
| 644 | line 23 | replace "Chapter 11" with "Chapter 10" | 14-9-00 |
| 670 | line -3 | replace "Seetharamanan" by "Seetharaman" | 19-7-02 |

| 673 | Section 17.2.1 | This example is described for Java 2 version 1.3 or earlier. See 'Java CORBA with Java 2 version 1.4' under 'Additional Material' on www.cdk3.net . (under Supplementary Material for Chapter 17.) | 2-4-03 |
|-----|-----------------------------|--|----------|
| 683 | Figure 17.8 lines -5 and -2 | add "{" and replace "strings" by "string s" | 28-11-00 |
| 698 | Exercise 17.19 | replace "17.13" by "Exercise 17.18" | 13-7-00 |
| 714 | line 14 | insert the following: "a 'copy-inherited' region is one that is" after "copied:" | 6-9-01 |
| 745 | Insert additional reference | Neumann and Ts'o 1994 Neuman, B.C. and Ts'o, T. 1994. Kerberos: An Authentication Service for Computer Networks, <i>IEEE</i> Communications, vol. 32, no. 9, pp. 33-38. Sept. 1994. http://nii.isi.edu/publications/kerberos-neuman-tso.html | 15-9-00 |
| 751 | penultimate reference | replace "Spasojevic R"by "Spasojevic M" | 2-8-00 |
| 751 | fifth reference | replace "Seetharamanan" by "Seetharaman" (twice) | 19-7-02 |

 ${\it Most recently reported error~14~April~2003~@George~Coulouris,~Jean~Dollimore~and~Tim~Kindberg~2000}$